

MORE SECRETS OF CLICKHOUSE QUERY PERFORMANCE

with Robert Hodges

Brief Intros



Robert Hodges - Altinity CEO

30+ years on DBMS plus virtualization and security.

ClickHouse is DBMS #20



www.altinity.com

Leading software and services provider for ClickHouse

Major committer and community sponsor in US and Western Europe



Goals of the talk

- Understand single node MergeTree structure
- Improve response by tuning queries
- Get much bigger gains by changing data layout
- Increase storage performance with new multi-disk volumes

Non-Goals:

- Boost performance of sharded/replicated clusters
- Teach advanced ClickHouse performance management



ClickHouse & MergeTree Intro





Introduction to ClickHouse

Understands SQL

Runs on bare metal to cloud

Shared nothing architecture

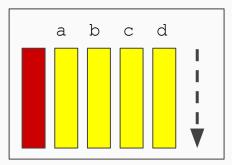
Stores data in columns

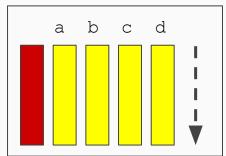
Parallel and vectorized execution

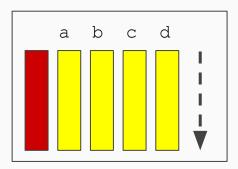
Scales to many petabytes

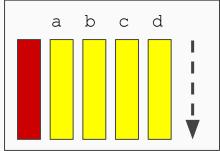
Is Open source (Apache 2.0)

And it's <u>really</u> fast!









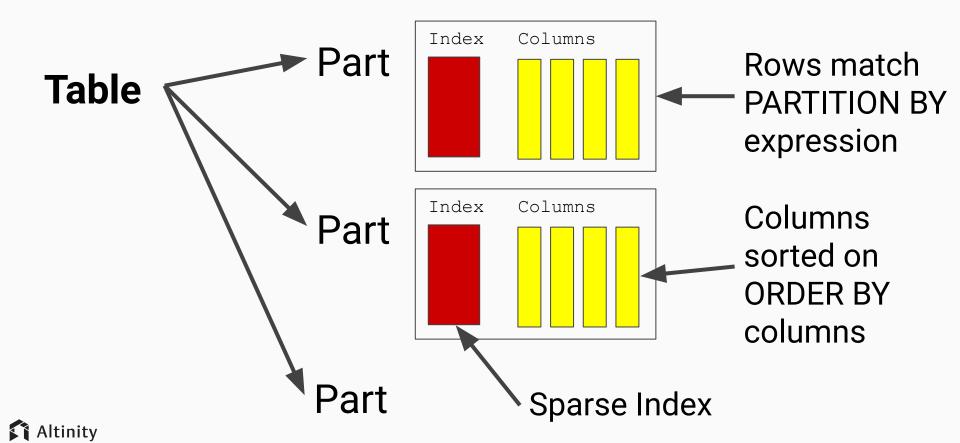


Introducing the MergeTree table engine

```
CREATE TABLE ontime (
  Year UInt16,
                                     Table engine type
  Quarter UInt8,
  Month UInt8,
                                     How to break data
                                     into parts
  ENGINE = MergeTree()
PARTITION BY to YYYYMM (FlightDate)
ORDER BY (Carrier, FlightDate)
                                     How to index and
                                     sort data in each part
```

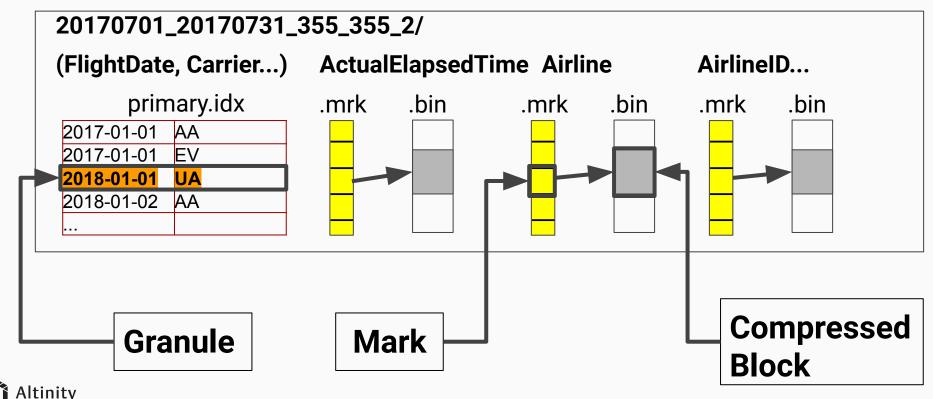


Basic MergeTree data layout



MergeTree layout within a single part

/var/lib/clickhouse/data/airline/ontime_reordered



Basic Query Tuning





ClickHouse performance tuning is different...

The bad news...

- No query optimizer
- No EXPLAIN PLAN
- May need to move [a lot of] data for performance

The good news...

- No query optimizer!
- System log is great
- System tables are too
- Performance drivers are simple: I/O and CPU
- Constantly improving



Your friend: the ClickHouse query log

Return messages to clickhouse-client clickhouse-client --send logs level=trace Must enable in select * from system.text log config.xml sudo less \ /var/log/clickhouse-server/clickhouse-server.log View all log messages on server



Use system log to find out query details

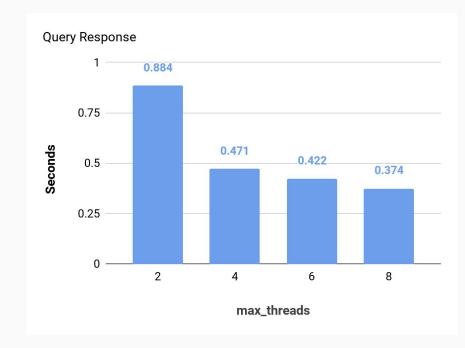
```
SELECT toYear(FlightDate) year,
  sum(Cancelled)/count(*) cancelled,
  sum(DepDel15)/count(*) delayed 15
FROM airline.ontime
GROUP BY year ORDER BY year LIMIT 10
                                         (Log messages)
 Query pipeline in log
                                         Limit
                                          Expression
                                           MergeSorting
                                            PartialSorting
                                             Expression
 8 parallel threads
                                              ParallelAggregating
 to read table
                                               Expression × 8
                                                MergeTreeThread
```



Speed up query executing by adding threads

```
SELECT toYear(FlightDate) year,
  sum(Cancelled)/count(*) cancelled,
  sum(DepDel15)/count(*) delayed 15
FROM airline.ontime
GROUP BY year ORDER BY year LIMIT 10
SET max threads = 2
SET max threads = 4
```

max_threads defaults to half the number of physical CPU cores





Speed up queries by reducing reads

```
SELECT toYear(FlightDate) year,
   sum(Cancelled)/count(*) cancelled,
   sum(DepDel15)/count(*) delayed_15
FROM airline.ontime
GROUP BY year ORDER BY year LIMIT 10
```

```
SELECT toYear(FlightDate) year,
   sum(Cancelled)/count(*) cancelled,
   sum(DepDel15)/count(*) delayed_15
FROM airline.ontime
WHERE year =
toYear(toDate('2016-01-01'))
GROUP BY year ORDER BY year LIMIT 10
```

```
(Log messages)
Selected 355 parts by date,
355 parts by key,
21393 marks to read from 355
ranges
```

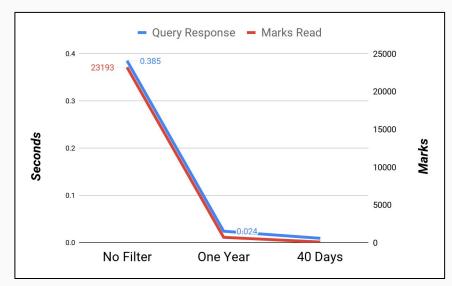
```
(Log messages)
Selected 12 parts by date,
12 parts by key,
692 marks to read from 12
ranges
```



Query execution tends to scale with I/O

```
SELECT
  FlightDate,
  count(*) AS total_flights,
  sum(Cancelled) / count(*) AS cancelled,
  sum(DepDel15) / count(*) AS delayed_15
FROM airline.ontime
WHERE (FlightDate >= toDate('2016-01-01'))
  AND (FlightDate <= toDate('2016-02-10'))
GROUP BY FlightDate</pre>
```

```
(Log messages)
Selected 2 parts by date,
2 parts by key,
73 marks to read from 2 ranges
```





Use PREWHERE to help filter unindexed data

```
Year, count(*) AS total_flights,
count(distinct Dest) as destinations,
count(distinct Carrier) as carriers,
sum(Cancelled) / count(*) AS cancelled,
sum(DepDel15) / count(*) AS delayed_15
FROM airline.ontime [PRE]WHERE Dest = 'BLI' GROUP BY Year
```

```
(PREWHERE Log messages)
Elapsed: 0.591 sec.
Processed 173.82 million rows,
2.09 GB (294.34 million rows/s.,
3.53 GB/s.)
```

```
(WHERE Log messages)
Elapsed: 0.816 sec.
Processed 173.82 million rows,
5.74 GB (213.03 million rows/s.,
7.03 GB/s.)
```



But PREWHERE can kick in automatically

```
SET optimize_move_to_prewhere = 1

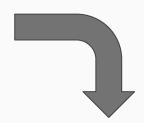
SELECT

Year, count(*) AS total_flights,
count(distinct Dest) as destinations,
count(distinct Carrier) as carriers,
sum(Cancelled) / count(*) AS cancelled,
sum(DepDel15) / count(*) AS delayed_15

FROM airline.ontime

WHERE Dest = 'BLI' GROUP BY Year
```

This is the default value



```
(Log messages)
InterpreterSelectQuery:
MergeTreeWhereOptimizer: condition
"Dest = 'BLI'" moved to PREWHERE
```

Restructure joins to reduce data scanning

```
SELECT
Dest d, Name n, count(*) c, avg(ArrDelayMinutes)
  FROM ontime
   JOIN airports ON (airports.IATA = ontime.Dest)
   GROUP BY d, n HAVING c > 100000 ORDER BY d DESC
     LIMIT 10
                        Faster
SELECT dest, Name n, c AS flights, ad FROM (
SELECT Dest dest, count(*) c, avg(ArrDelayMinutes) ad
 FROM ontime
  GROUP BY dest HAVING c > 100000
    ORDER BY ad DESC
) LEFT JOIN airports ON airports.IATA = dest LIMIT 10
```

3.878 seconds

1.177 seconds



The log tells the story

Join <u>during</u> MergeTree scan



Join <u>after</u> MergeTree scan



```
(Log messages)
ParallelAggregatingBlockInputStream
: Total aggregated. 173790727 rows
(from 10199.035 MiB) in 3.844 sec.
(45214666.568 rows/sec., 2653.455
MiB/sec.)
```

(Log messages)
ParallelAggregatingBlockInputStream
: Total aggregated. 173815409 rows
(from 2652.213 MiB) in 1.142 sec.
(152149486.717 rows/sec., 2321.617
MiB/sec.)



More ways to find out about queries

```
Start query logging
SET log queries = 1
Run a query
SELECT version()
                                          Stop query logging
SET log queries = 0
SELECT * FROM system.query log
  WHERE query='SELECT version()'
                                          Show currently
SHOW PROCESSLIST
                                          executing queries
```



Optimizing Data Layout





Restructure data for big performance gains

- Ensure optimal number of parts
- Optimize primary key index and ordering to reduce data size and index selectivity
- Use skip indexes to avoid unnecessary I/O
- Use encodings to reduce data size before compression
- Use materialized views to transform data outside of source table
- Plus many other tricks



How do partition keys affect performance?

```
CREATE TABLE ontime ...
ENGINE=MergeTree()
PARTITION BY
  toYYYYMM(FlightDate)
                                                  Is there a
                                                   practical
                                                  difference?
CREATE TABLE ontime many parts
ENGINE=MergeTree()
PARTITION BY FlightDate
```



Keep parts in the hundreds, not thousands

CREATE TABLE ontime ...

ENGINE=MergeTree()

PARTITION BY

toYYYYMM(FlightDate)



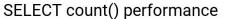
CREATE TABLE ontime_many_parts

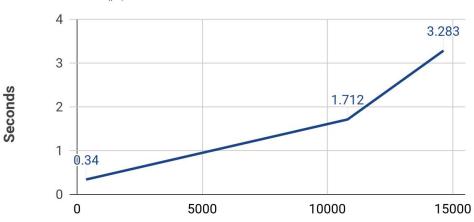
. . .

ENGINE=MergeTree()

PARTITION BY FlightDate







Number of Parts

Table	Rows	Parts
ontime	174M	355
ontime_many_parts (after OPTIMIZE)	174M	10,085
ontime_many_parts (before OPTIMIZE)	174M	14,635

Think about primary key index structure

```
CREATE TABLE ontime_reordered (
   Year UInt16,
   Quarter` UInt8,
   . . .)

ENGINE = MergeTree()

PARTITION BY toYYYYMM(FlightDate)

ORDER BY (Carrier, Origin, FlightDate)

SETTINGS index_granularity=8192
```

Hashing large values allows index to fit in memory more easily

Small granularity can make skip indexes more selective

Large granularity makes index smaller



Table ORDER BY is key to performance

```
CREATE TABLE ontime_reordered (
   Year UInt16,
   Quarter` UInt8,
   . . .)
ENGINE = MergeTree()
PARTITION BY toYYYYMM(FlightDate)
ORDER BY (Carrier, Origin, FlightDate)
SETTINGS index_granularity=8192
```

Choose order to make dependent non-key values less random

Benefits:

- → Higher compression
- → Better index selectivity
- → Better PREWHERE performance



Skip indexes cut down on I/O

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```
SET allow experimental data skipping indices=1;
                                                     Default value
ALTER TABLE ontime ADD INDEX
  dest name Dest TYPE ngrambf v1(3, 512, 2, 0) GRANULARITY 1
ALTER TABLE ontime ADD INDEX
  cname Carrier TYPE set(0) GRANULARITY 1
OPTIMIZE TABLE ontime FINAL
-- OR, in current releases
ALTER TABLE ontime
  UPDATE Dest=Dest, Carrier=Carrier
    WHERE 1=1
```

Indexes & PREWHERE remove granules

Apply PREWHERE on Dest predicate

```
(Log messages)
InterpreterSelectQuery: MergeTreeWhereOptimizer:
condition "Dest = 'PPG'" moved to PREWHERE
. . . .
(SelectExecutor): Index `dest_name` has dropped 55
granules.
(SelectExecutor): Index `dest_name` has dropped 52
granules.
```

Use index to remove granules from scan



Effectiveness depends on data distribution

```
Year, count(*) AS flights,
sum(Cancelled) / flights AS cancelled,
sum(DepDel15) / flights AS delayed_15
FROM airline.ontime WHERE [Column] = [Value] GROUP BY Year
```

Column	Value	Index	Count	Rows Processed	Query Response
Dest	PPG	ngrambf_v1	525	4.30M	0.053
Dest	ATL	ngrambf_v1	9,360,581	166.81M	0.622
Carrier	ML	set	70,622	3.39M	0.090
Carrier	WN	set	25,918,402	166.24M	0.566



Current index types

Name	What it tracks
minmax	High and low range of data; good for numbers with strong locality like timestamps
set	Unique values
ngrambf_v1	Presence of character ngrams, works with =, LIKE, search predicates; good for long strings
tokenbf_v1	Like ngram but for whitespace-separated strings; good for searching on tags
bloomfilter	Presence of value in column

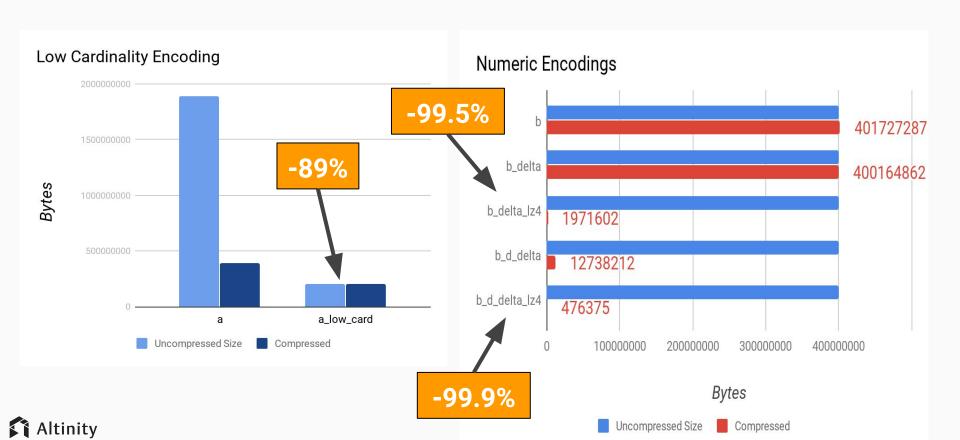


Encodings improve compression efficiency

```
Values with
CREATE TABLE test codecs (a String,
                                                        dictionary
   a lc LowCardinality(String) DEFAULT a,
                                                        encoding
   b UInt32,
   b delta UInt32 DEFAULT b Codec(Delta),
   b delta 1z4 UInt32 DEFAULT b Codec (Delta, LZ4),
                                                        Differences
   b dd UInt32 DEFAULT b Codec (DoubleDelta),
                                                        between
   b dd 1z4 UInt32 DEFAULT b Codec (DoubleDelta, LZ4)
                                                        values
Engine = MergeTree
PARTITION BY tuple() ORDER BY tuple();
                                                   Differences
                                                   between change
                                                   of value
```

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Effect on storage size is dramatic



Queries are faster due to less I/O

```
SELECT a AS a, count(*) AS c FROM test_codecs
GROUP BY a ORDER BY c ASC LIMIT 10
. . .
```

10 rows in set. Elapsed: 0.681 sec. Processed 100.00 million rows, 2.69 GB (146.81 million rows/s., 3.95 GB/s.)



```
SELECT a_lc AS a, count(*) AS c FROM test_codecs
GROUP BY a ORDER BY c ASC LIMIT 10
```

10 rows in set. Elapsed: 0.148 sec. Processed 100.00 million rows, 241.16 MB (675.55 million rows/s., 1.63 GB/s.)



Overview of encodings

Name	Best for
LowCardinality	Strings with fewer than 10K values
Delta	Time series
Double Delta	Increasing counters
Gorilla	Gauge data (bounces around mean)
T64	Integers other than random hashes



Compression may vary across ZSTD and LZ4

TIP: use system.columns to check data size

```
SELECT table,
  formatReadableSize(sum(data_compressed_bytes)) tc,
  formatReadableSize(sum(data_uncompressed_bytes)) tu,
  sum(data_compressed_bytes) / sum(data_uncompressed_bytes) as ratio
FROM system.columns
WHERE database = currentDatabase()
GROUP BY table ORDER BY table
```



Use mat views to boost performance further

```
CREATE MATERIALIZED VIEW ontime_daily_cancelled_mv

ENGINE = SummingMergeTree

PARTITION BY tuple() ORDER BY (FlightDate, Carrier)

POPULATE

AS SELECT

FlightDate, Carrier, count(*) AS flights,
sum(Cancelled) / count(*) AS cancelled,
sum(DepDel15) / count(*) AS delayed_15

FROM ontime

GROUP BY FlightDate, Carrier
```

Returns cancelled/late flights where Carrier = 'WN' in 0.007 seconds



More things to think about

Use smaller datatypes wherever possible

Use ZSTD compression (slower but better ratio)

Use dictionaries instead of joins

Use sampling when approximate answers are acceptable

Shard/replicate data across a cluster for large datasets

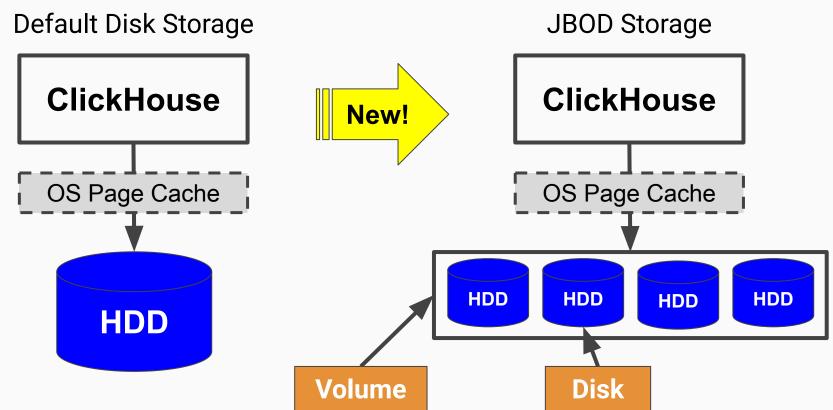
Check out "Further Resources" slide for more information



Thinking about Storage and Memory



ClickHouse now has flexible storage policies





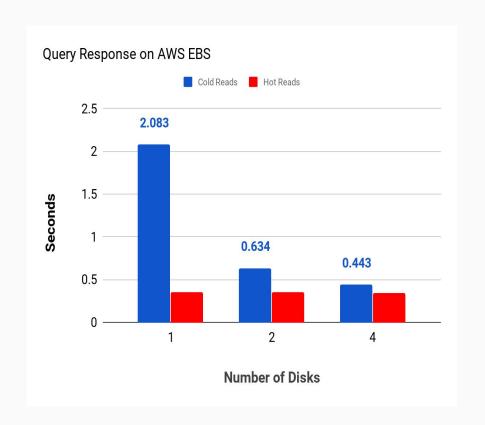
How do you apply storage policies?

```
CREATE TABLE. tripdata
  `pickup date` Date DEFAULT toDate(tpep pickup datetime),
ENGINE = MergeTree()
PARTITION BY to YYYYMM (pickup date)
ORDER BY (pickup location id, dropoff location id, vendor id)
SETTINGS storage policy = 'ebs jbod 4', -
         index granularity = 8192
                                                          Policy
```



Effect of storage policies on response

```
-- Cold query
set min bytes to use direct io=1
SELECT avg(passenger count)
FROM tripdata
  Hot query
set min bytes to use direct io=0
SELECT avg(passenger count)
FROM tripdata
SELECT avg (passenger count)
FROM tripdata
```



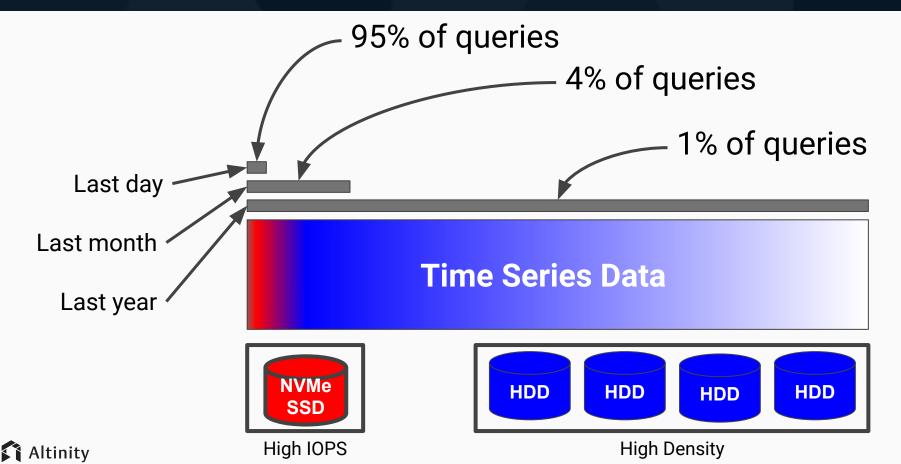


TIP: system.parts tracks content across disks

```
SELECT
  database, table, disk name,
  count(*) AS parts,
  uniq(partition) AS partitions,
  sum(marks) AS marks,
  sum(rows) AS rows,
  formatReadableSize(sum(data compressed bytes)) AS compressed,
  formatReadableSize(sum(data uncompressed bytes)) AS uncompressed,
  round(sum(data compressed bytes) / sum(data uncompressed bytes) * 100.0, 2)
AS percentage
FROM system.parts
WHERE active and database = currentDatabase()
GROUP BY database, table, disk name
ORDER BY database ASC, table ASC, disk name ASC
```



Tiered storage is another new option



And don't forget all the great OS utilities!

- top and htop -- CPU and memory
- dstat -- I/O and network consumption
- iostat I/O by device
- iotop -- I/O by process
- iftop -- Network consumption by host
- perf top -- CPU utilization by system function

For a full description see <u>Performance Analysis of ClickHouse Queries</u> by Alexey Milovidov



Wrap-up



Takeaways on ClickHouse Performance

- ClickHouse performance drivers are CPU and I/O
- The system query log is key to understanding performance
- Query optimization can improve response substantially
- Restructure tables and add indexes/mat views for biggest gains
- In recent versions you can now optimize storage, too!



Further resources

- Altinity Blog, especially:
 - Amplifying ClickHouse Capacity with Multi-Volume Storage
- Altinity Webinars, especially:
 - ClickHouse Materialized Views: The Magic Continues
 - Strength in Numbers: Introduction to ClickHouse Cluster Performance
- ClickHouse documentation
- <u>Performance Analysis of ClickHouse Queries</u> by Alexey Milovidov
- ClickHouse Telegram Channel
- ClickHouse Slack Channel



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